

# Networks - Week 6 - Discrete Random Walks on Networks

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November 2023

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# 1 Formalising Discrete Random Walks

## 1.1 Discrete Random Walks on Networks

- Given a transition matrix, how can we compute the probability of a random walker arriving at a given vertex?

- let  $T$  be a **transition matrix**, such that:

$$T_{ij} = \text{probability that the walker goes from } i \text{ to } j$$

- the **probability** of the walker reaching the  $i$ th node in  $t$  steps is  $p_i(t)$ , which is the component of the matrix:

$$\underline{p}(t) = \underline{p}(0)T^t$$

- this section focuses on understanding  $\underline{p}$  through the **spectral properties** of  $T$

- What is the transition matrix for a random walk?

- given a **random walk**, where the walker will uniformly randomly pick an adjacent edge:

$$T_{ij} = \frac{A_{ij}}{d_i}$$

where  $d_i$  denotes the **degree** of the  $i$ th vertex

- When is the transition matrix symmetric?

- whenever the underlying network is **symmetric** (need an undirected graph) and whenever each vertex has the same degree

- for  $T_{ij} = T_{ji}$  we need that:

$$A_{ij} = A_{ji} \quad d_i = d_j$$

- If the transition matrix is not symmetric, how can we derive its spectral properties?

- **symmetric matrices** have very nice, well-studied **spectral properties** (**Perron-Frobenius Theorem**, guarantee on real eigenvalues, etc...)

- even if  $T$  isn't symmetric, we can still study its properties through a **related** matrix which is symmetric:

$$\tilde{A}_{ij} = \frac{A_{ij}}{\sqrt{d_i d_j}}$$

(whenever  $A$  is a symmetric matrix)

- notice, the **spectral properties** of  $\tilde{A}$  will be essentially equivalent to those of the **normalised Laplacian**

Let  $A$  be an **adjacency matrix**. The **Laplacian** of  $A$  is constructed by defining a matrix

$$D = \text{diag}(d_1, \dots, d_n)$$

where  $d_i$  is the  $i$ th row sum of  $A$ . The Laplacian of  $A$  is the **symmetric** matrix:

$$L = D - A$$

The **normalised Laplacian** is given by:

$$\tilde{L} = I - D^{-1/2} A D^{-1/2}$$

where since  $D$  is a diagonal matrix:

$$D^\omega = \text{diag}(d_1^\omega, \dots, d_n^\omega)$$

where we identify:

$$\tilde{A} = D^{-1/2} A D^{-1/2}$$

- How can we express  $\tilde{A}$  in terms of its spectrum?

- since  $\tilde{A}$  is **real** and **symmetric**, we can **diagonalise** it:

$$\tilde{A} = U \Lambda U^T$$

- if  $\{\lambda_i\}_{i \in [1, n]}$  gives a set of **eigenvalues**, and  $\{\underline{u}_i\}_{i \in [1, n]}$  gives a set of **eigenvectors**, then the **orthonormal decomposition** of  $\tilde{A}$  is:

$$\tilde{A} = \sum_{i=1}^n \lambda_i \underline{u}_i \underline{u}_i^T$$

- How can we use the spectrum of  $\tilde{A}$  to express  $T$ ?

- notice:

$$T_{ij} = \frac{\sqrt{d_j} \tilde{A}_{ij}}{\sqrt{d_i}}$$

so in particular, if  $D = \text{diag}(d_1, \dots, d_n)$  denotes the diagonal matrix of degrees of  $A$ :

$$T = D^{-\frac{1}{2}} \tilde{A} D^{\frac{1}{2}}$$

- in particular, the **eigenvalues** of  $T$  are the same as those for  $\tilde{A}$  and:

- \* the **right eigenvectors** of  $T$  are:

$$\underline{w}_i^R = D^{-\frac{1}{2}} \underline{u}_i$$

since:

$$D^{-\frac{1}{2}} \tilde{A} D^{\frac{1}{2}} (D^{-\frac{1}{2}} \underline{u}_i) = D^{-\frac{1}{2}} (\tilde{A} \underline{u}_i) = \lambda_i D^{-\frac{1}{2}} \underline{u}_i$$

- \* similarly, the **left eigenvectors** of  $T$  are:

$$\underline{w}_i^L = D^{\frac{1}{2}} \underline{u}_i$$

- this gives us the **eigendecomposition**

$$T = \sum_{i=1}^n \lambda_i \underline{w}_i^R (\underline{w}_i^L)^T$$

- Using the eigendecomposition of the transition, how may we express the walker probability matrix  $\underline{p}(t)$  (known as the stationary density)?

- using the **eigendecomposition**, we have that:

$$T^t = \sum_{i=1}^n \lambda_i^t \underline{w}_i^R (\underline{w}_i^L)^T$$

- hence, if  $\underline{p}(t)$  is a **row vector**:

$$\underline{p}(t) = \underline{p}(0) T^t = \sum_{i=1}^n \lambda_i^t (\underline{w}_i^L)^T (\underline{p}(0) \underline{w}_i^R)$$

- In practice, how is the stationary density computed?

- via a **power method**
- that is, start with a **random vector**  $\underline{v}_0$
- then define:

$$\underline{v}_{i+1} = T \underline{v}_i$$

## 1.2 Proposition: Spectral Information of the Transition Matrix

Let  $T$  be the **transition matrix** given by:

$$T_{ij} = \frac{A_{ij}}{d_i}$$

Then:

1. All **eigenvalues**  $\lambda_i$  of  $T$  lie in the interval  $[-1, 1]$
2. The **left eigenvector**  $(\underline{w}_j^L)^T$  corresponding to the **eigenvalue**

$$\lambda_j = 1$$

gives the **stationary density** of the **transition matrix**; that is we have:

$$(\underline{w}_j^L)^T = \underline{p}^*$$

whereby:

$$\lim_{t \rightarrow \infty} \underline{p}(t) = \underline{p}^* \quad \text{and} \quad \underline{p}^* = \underline{p}^* T$$

3. The **right eigenvector**  $\underline{w}_j^R$  corresponding to  $\lambda_j = 1$  is such that:

$$\underline{w}_j^R \propto \begin{pmatrix} 1 \\ \vdots \\ 1 \end{pmatrix}$$

4. The **eigenvalue**  $\lambda_j$  is the **largest magnitude eigenvalue**, and all elements of  $\underline{w}_j^R, \underline{w}_j^L$  are **positive**.

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*Proof.* We make extensive use of the Perron-Frobenius Theorem, since we assume that  $T$  is irreducible (i.e. that there is always a path between any 2 vertices  $i, j$ , or that  $\exists n : T^n = 0$ ).

Let  $A$  be a  $n \times n$  matrix, such that  $A$ :

- is **irreducible**
- **non-negative**
- has **spectral radius**  $\rho(A) = r > 0$

Then:

1.  $r$  is an **eigenvalue** of  $A$  (called the **Perron-Frobenius eigenvalue**)
2.  $r$  is **simple**. In particular:
  - $r$  has **algebraic multiplicity 1** (it is not a repeated eigenvalue)
  - $r$  has **geometric multiplicity 1** (both right and left eigenspaces are one-dimensional - this is because geometric multiplicity is bounded by algebraic multiplicity)
3.  $A$  has left/right **eigenvectors** with **eigenvalue**  $r$ , and whose components are **all positive**
4. the only **eigenvectors** whose components are **all positive** are those associated to  $r$
5.  $r$  is **bounded** above/below by the maximum and minimum **row sums** of  $A$  (and also the **column sums**):

$$\min_{i \in [1, n]} \sum_{j=1}^n A_{ij} \leq r \leq \max_{i \in [1, n]} \sum_{j=1}^n A_{ij}$$

①

By Perron-Frobenius (part 5), the Perron-Frobenius Eigenvalue of  $T$  is bounded by row/column sums of  $T$ . In particular:

$$r \leq \sum_{j=1}^n \frac{A_{ij}}{d_i} = \frac{1}{d_i} \sum_{j=1}^n A_{ij} \frac{d_i}{d_i} = 1$$

where we've used the definition of degree as:

$$\sum_{j=1}^n \frac{A_{ij}}{d_i} d_i$$

In particular, since the eigenvalues of  $T$  are all real (as they are eigenvalues of a real symmetric matrix  $\tilde{A}$ ), and the absolute value of all eigenvalues is bounded by the spectral radius (which is given by the Perron-Frobenius Eigenvalue), it follows that any eigenvalue of  $T$  satisfies:

$$\lambda_i \in [-1, 1]$$

②

By definition of the stationary density:

The **non-negative stationary density** is a vector:

$$\underline{p}^* = (p_1^*, \dots, p_n^*)$$

where:

$$p_i^* = \lim_{t \rightarrow \infty} p_i(t)$$

and:

$$\underline{p}^* = \underline{p}^* T$$

it corresponds to the left eigenvector of the transition matrix with eigenvalue 1. This is precisely  $(\underline{w}_j^L)^T$ .

③

If  $\underline{w}_j^R$  is the right eigenvector of  $T$  with eigenvalue 1, we have that:

$$T \underline{w}_j^R = \underline{w}_j^R$$

Notice also that if

$$\underline{v} = c \begin{pmatrix} 1 \\ \vdots \\ 1 \end{pmatrix}$$

Then:

$$T(\underline{c}\underline{v}) = \underline{c}\underline{v}$$

since the row sums of  $T$  are 1 (as each row element is a probability). Hence,  $\underline{c}\underline{v}$  is a right eigenvector of  $T$  with eigenvalue 1, so in particular:

$$\forall c \in \mathbb{R}, \quad \underline{w}_j^R = \underline{c}\underline{v}$$

④

This is properties 1 and 3 of the Perron-Frobenius Theorem. □

## 1.3 The Cheeger Constant and the Spectral Gap

### 1.3.1 The Spectral Gap

- How can we approximate the stationary density for large  $t$ ?

– by the **Perron-Frobenius Theorem**, the **eigenvalues** of  $T$  are bounded on:

$$\lambda_i \in [-1, 1]$$

– in particular, since the **stationary density** is given by:

$$\underline{p}(t) = \sum_{i=1}^n \lambda_i^t (\underline{w}_i^L)^T (\underline{p}(0) \underline{w}_i^R)$$

as  $t \rightarrow \infty$ , most eigenvalues will tend to 0

– thus, if  $\lambda_{max} = 1$ , and  $\lambda_2$  is the second largest eigenvalue (in magnitude), it follows that:

$$\underline{p}(t) \approx (\underline{w}_{max}^L)^T (\underline{p}(0) \underline{w}_{max}^R) + \lambda_2^t (\underline{w}_2^L)^T (\underline{p}(0) \underline{w}_2^R)$$

- What is the spectral gap?

- the value:

$$\gamma = \lambda_{max} - |\lambda_2| = 1 - |\lambda_2|$$

- What does a large spectral gap imply?

- if  $\gamma$  is large, the term:

$$\lambda_2^t (\underline{w}_2^L)^T (\underline{p}(0) \underline{w}_2^R)$$

goes to 0 much quicker (since  $|\lambda_2|$  is small)

- in particular,  $\lambda_2$  controls the “relaxation time” (i.e the time required for the stationary density to stabilise into a constant)

### 1.3.2 Definition: The Cheeger Inequality

The Cheeger Inequality gives useful bounds on  $|\lambda_2|$ , and thus informs us on the relaxation time.

Let  $S$  be a **vertex subset** of a **network**, and let  $\bar{S}$  be its **complement**. The **Cheeger Constant** (also known as **conductance**) is defined by:

$$h = \min_S \left\{ \frac{\text{of edges connecting } S \text{ and } \bar{S}}{\min\{|S|, |\bar{S}|\}} \right\}$$

The **Cheeger Inequality** is:

$$\frac{h^2}{2} < 1 - |\lambda_2| \leq 2h$$

---

*Intuitively, this makes sense. If  $h$  is small, this implies that there are a small number of edges connecting  $S, \bar{S}$ . This can be interpreted as the graph being split into 2 communities (namely  $S, \bar{S}$ ) with a few edges connecting them. In such a scenario, the probability that a random walker moves between communities is low, which means that so too the spectral gap will be small, so in particular the relaxation time will be longer.*

## 2 Applications of Random Walks

### 2.1 The PageRank Algorithm

#### 2.1.1 Motivating PageRank

- What is the purpose of the PageRank algorithm?

- define a **centrality measure** on **directed networks**
- **PageRank** is defined such that a node  $v_i$  is “important” if:
  - \* it has many **incoming edges**
  - \* the **incoming edges** originate from **important nodes**



- \* if  $v_j$  has an **outgoing edge** to  $v_i$ , the **out-degree** of  $v_j$  is small
- in short,  $v_i$  has to be connected to a lot of important nodes, which themselves don't share many connections with other nodes
- these requirements were used to **rank** webpages, based on the **hyperlinks** which connected the websites
- **What is the PageRank of a given node?**
  - the corresponding **stationary density** of said node, given a **discrete-time random walk** through the **directed network**
- **Why can't we immediately derive a stationary density  $\underline{p}^*$  in the case of PageRank (i.e in some random directed network)?**
  - the **stationary density** exists and is **unique** if and only if the network is **strongly connected**
  - otherwise:
    - \* many possible  $\underline{p}^*$  can exist (i.e if the network contains multiple absorbing states, whereby  $T_{ii} = 1$ )
    - \* certain  $p_i^*$  will be 0 (i.e they will be considered as irrelevant)
  - generally, **real world networks** (and especially the Internet) don't satisfy **strong connectedness**

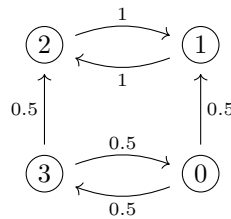


Figure 1: Example of the **spider trap** problem, whereby 2 nodes will get  $p^* = 0$ . This is because once states 1, 2 are entered, they will indefinitely cycle amongst themselves. In particular, this is an example of a **periodic Markov Chain**, and it yields  $\underline{p}^* = \langle 0, 0.5, 0.5, 0 \rangle$ ; here, two states take away all the “importance” of the other two states.

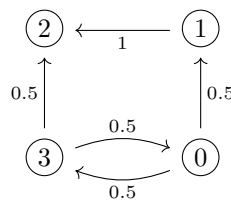


Figure 2: Example of the **dead ends** problem, whereby node 2 is an absorbing state. In such a case, the algorithm will fail to converge and no unique stationary density vector will be found. This is because once state 2 is reached, it can never be escaped; thus, different  $\underline{p}^*$  can be found through **power iteration** methods, depending on the initial conditions of the process.

- **How can lack of strong connectivity be tackled?**
  - **PageRank** can “ignore” the lack of **strong connectivity** by assuming **node teleportation**: from a given website, you can randomly “jump” to any other website (i.e you don't use a hyperlink to find a new website, you just search it up)
  - this reasonable assumption then allows us to find **unique stationary density**

### 2.1.2 Definition: Random Walk with Teleportation

Suppose we have a **directed network** traversed by a **random walker** who's dynamics are defined by the **transition matrix**:

$$T_{ij} = \frac{A_{ij}}{d_i^{\text{out}}}$$

Then, **random walks with teleportation** are defined by:

$$\underline{p}(t+1) = \alpha \underline{p}(t)T + (1-\alpha)\underline{b}$$

or componentwise:

$$p_i(t+1) = \alpha \sum_{j=1}^n p_j(t)T_{ji} + (1-\alpha)b_i$$

where:

- $\alpha \in (0, 1)$  is the **damping factor**, which is the probability that at any given time-step the walker follows links (and  $1 - \alpha$  is thus the probability that it teleports)
- $\underline{b}$  is a **preference row vector** satisfying:

$$\sum_{i=1}^n b_i = 1$$

and such that  $b_i \in [0, 1]$  is the probability of randomly teleporting to the  $i$ th vertex

- 
- What value is typically used for the damping factor?

- a well-agreed value is  $\alpha = 0.85$

### 2.1.3 Proposition: PageRank Stationary State

The **PageRank Stationary State** can be calculated via:

$$\underline{p}^* = (1 - \alpha)\underline{b}(I - \alpha T)^{-1}$$

or componentwise:

$$p_i^* = b_i + \sum_{l=1}^{\infty} \sum_{j=1}^n b_j (T_{ji} - T_{ji}^{l-1})$$

---

Notice, here we assume that  $\alpha$  is well chosen, such that the geometric series form of  $(I - \alpha T)^{-1}$  is well-defined, like in Katz Centrality

## 2.2 Epidemic Modelling

### 2.2.1 The SIS Model

#### 2.2.1.1 Definition: SIS Model

Consider a **clique** of  $n$  vertices. Let:

- $S(t)$  be the **proportion** of nodes which are **susceptible** (can be infected) at time  $t$
- $I(t)$  be the **proportion** of nodes which are **infected** (can infect) at time  $t$

The **dynamics** of the **SIS Model** (susceptible-infected-susceptible) are defined by a system of ODEs:

$$\begin{aligned}\frac{dS}{dt} &= -\beta I(t)S(t) + \mu I(t) \\ \frac{dI}{dt} &= \beta I(t)S(t) - \mu I(t)\end{aligned}$$

where:

- $\beta$  is the **infection rate** (the rate at which a **susceptible node** contracts an infection, and thus becomes an **infected node**)
- $\mu$  is the **recovery rate** (the rate at which an **infected node** recovers from the infection, and thus becomes a **susceptible node**)

### 2.2.1.2 Explaining the SIS Model

- What is the probability of infection on some small time frame  $\Delta t$ ?

– we can approximate this probability via:

$$\beta \Delta t$$

- What is the intuition behind the ODEs defining the SIS model?

– for the **susceptible nodes**:

- \* if  $S$  increases, then we expect the rate of change of  $S$  to decrease (since more susceptible nodes are available to infect)
- \* increasing  $I$  will both **decrease** the number of susceptibles (since more nodes can now infect at a rate  $\beta$ ) and **increase** the number of susceptibles (since more nodes can now recover at a rate  $\mu$ )
- \* thus, if  $\beta \gg \mu$ , the number of susceptibles will generally decrease, but if  $\mu \gg \beta$ , the number of susceptibles will generally increase

– similar arguments work for the **infected nodes**, but the other way round

– in particular, we have that:

$$\frac{d}{dt}(S + I) = 0 \implies \forall t \geq 0, S + I = 1$$

(since  $S(0) + I(0) = 1$ )

- What is equilibrium like for the SIS model?

– let  $S^*, I^*$  be equilibrium values (that is, points at which  $S'(t) = I'(t) = 0$ ).

– then, assuming that  $I^* \neq 0$ :

$$\begin{aligned} 0 &= -\beta I^* S^* + \mu I^* & 0 &= \beta I^* S^* - \mu I^* \\ \implies \beta S^* &= \mu \\ \implies S^* &= \frac{\mu}{\beta} \end{aligned}$$

and since  $S + I = 1$  always, we must then have that:

$$S^* = \frac{\mu}{\beta} \quad I^* = 1 - \frac{\mu}{\beta}$$

– otherwise, we will have that  $I^* = 0$  (i.e at equilibrium there are no infections)

- What is the epidemic threshold of the SIS model?

– the **epidemic threshold**  $\varepsilon$  is a value such that if  $\varepsilon > 1$ , an **epidemic** is possible; if  $\varepsilon < 1$ , then sustaining the infections becomes impossible

– in the **SIS model**, we have that:

$$\varepsilon = \frac{\beta}{\mu}$$

since if:

$$\frac{\beta}{\mu} > 1$$

then  $S^*, I^*$  remain in the range  $(0, 1)$  and we have epidemic activity

– otherwise, we see that  $S^* > 1, I^* < 0$  and the epidemic is unsustainable

- Does the epidemic threshold make sense?

- suppose that

$$\frac{\beta}{\mu} > 1$$

such that the **epidemic** holds

- if  $\mu$  is **large** (i.e high recovery rate) or  $\beta$  is **small** (i.e low infection rate), we see that:

$$S^* \rightarrow 1 \quad I^* \rightarrow 0$$

- conversely, for **small**  $\mu$  or **large**  $\beta$ :

$$S^* \rightarrow 0 \quad I^* \rightarrow 1$$

## 2.2.2 The SIR Model

### 2.2.2.1 Definition: SIR Model

Consider a **clique** of  $n$  vertices. Let:

- $S(t)$  be the **proportion** of nodes which are **susceptible** (can be infected) at time  $t$
- $I(t)$  be the **proportion** of nodes which are **infected** (can infect) at time  $t$
- $R(t)$  be the **proportion** of nodes which are **recovered** (were infected, recovered and can no longer infect or be infected) at time  $t$

The **dynamics** of the **SIR Model** (susceptible-infected-recovered) are defined by a system of ODEs:

$$\begin{aligned}\frac{dS}{dt} &= -\beta I(t)S(t) + \mu I(t) \\ \frac{dI}{dt} &= \beta I(t)S(t) - \mu I(t) \\ \frac{dR}{dt} &= \mu I(t)\end{aligned}$$

where:

- $\beta$  is the **infection rate** (the rate at which a **susceptible node** contracts an infection, and thus becomes an **infected node**)
- $\mu$  is the **recovery rate** (the rate at which an **infected node** recovers from the infection, and thus becomes a **susceptible node**)

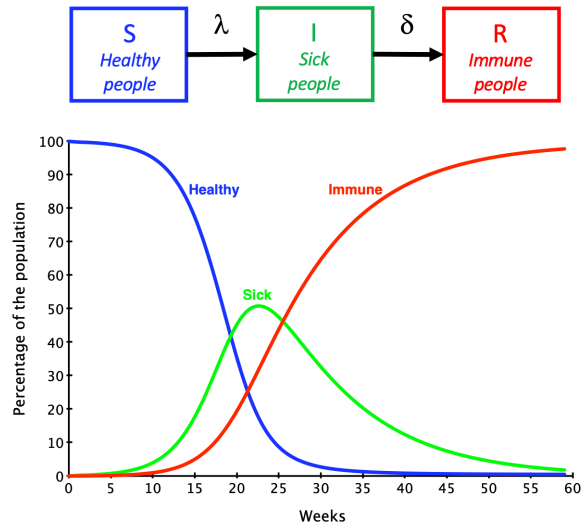


Figure 3: Dynamics of the **SIR Model**. Notice, the **infections** always reach a peak of contagion. The idea behind **epidemic management** is to spread the peak as much as possible, to avoid strain on medical services.

#### 2.2.2.2 Explaining the SIR Model

- What is the key dynamic difference between the SIS and SIR models?

- in the **SIR model**, the infection eventually **dies out**, even if  $\beta$  is large
- the final state solely consists of **susceptible** and **recovered** nodes

- When does the peak observed above occur?

- if

$$\frac{dI}{dt} > 0$$

then there will be a **perceptible** infectious outbreak

- otherwise, the rate of infection won't lead to outbreaks on a visible scale